## Transmathematics

and how to execute a program in less than one clock cycle Dr James Anderson FBCS CITP

## Agenda

- What is transmathematics?
- Acceptance of the complex numbers
- Acceptance of the transreal numbers
- Execute a program in less than one clock cycle


## Transmathematics

- Deals with total systems that have no exceptions
- Dividing by zero is not an exception
- Solves physical problems at singularities
- Allows exception-free programming
- Enables unbreakable pipelines


## Transmathematics

- Transmathematica journal
- Transmathematica conference
- Totallica company


## Complex Numbers

- It used to be impossible to find the square root of a negative number
- $\sqrt{-4}=$ ?


## Complex Numbers

- Define $i=j=\sqrt{-1}$


## Complex Numbers

- Now $(i 2)^{2}=i^{2} 2^{2}=-1 \times 4=-4$
- So $\sqrt{-4}=i 2$
- But people argued about this for 400 years!


## Complex Acceptance

- Complex numbers were constructed from real numbers
- So complex arithmetic is consistent if real arithmetic is
- So complex arithmetic cannot be disproved in its own terms, only real arithmetic can be disproved


## Complex Acceptance

- Controversy over complex numbers ended when they were given a geometrical interpretation


## Complex Geometry



## Complex Acceptance

- Fully accepted when useful applications were developed


## Complex Acceptance

- Construction from the real numbers
- Geometrical models
- Today, machine proofs of consistency
- Today, many useful applications


## Transreal Numbers

Transreal numbers, t , are proper fractions of real numbers, with a non-negative denominator, d, and a numerator, $n$, that is one of $-1,0,1$ when $d=0$

$$
t=\frac{n}{d}
$$

With k a positive constant:

$$
-\infty=\frac{-k}{0}=\frac{-1}{0}
$$

$$
\Phi=\frac{0}{0}
$$

$$
\infty=\frac{k}{0}=\frac{1}{0}
$$

## Negative Denominators

An improper fraction may have a negative denominator ( $-k$ ) which must be made positive before any transarithmetical operator is applied

$$
\frac{n}{-k}=\frac{-n}{-(-k)}=\frac{-1 \times n}{-1 \times(-k)}=\frac{-n}{k}
$$

# Multiplication 

$\frac{a}{b} \times \frac{c}{d}=\frac{a c}{b d}$

## Division

$$
\frac{a}{b} \div \frac{c}{d}=\frac{a}{b} \times \frac{d}{c}
$$

## Addition of Two Infinities

$$
\begin{gathered}
\infty+\infty=\frac{1}{0}+\frac{1}{0}=\frac{1+1}{0}=\frac{2}{0}=\frac{1}{0}=\infty \\
\infty+(-\infty)=\frac{1}{0}+\frac{-1}{0}=\frac{1-1}{0}=\frac{0}{0}=\Phi \\
-\infty+\infty=\frac{-1}{0}+\frac{1}{0}=\frac{-1+1}{0}=\frac{0}{0}=\Phi \\
-\infty+(-\infty)=\frac{-1}{0}+\frac{-1}{0}=\frac{-1+(-1)}{0}=\frac{-2}{0}=\frac{-1}{0}=-\infty
\end{gathered}
$$

## General Addition

$$
\frac{a}{b}+\frac{c}{d}=\frac{a d+b c}{b d}
$$

## Subtraction

$$
\frac{a}{b}-\frac{c}{d}=\frac{a}{b}+\frac{-c}{d}
$$

## Associativity

$$
\begin{aligned}
& a+(b+c)=(a+b)+c \\
& a \times(b \times c)=(a \times b) \times c
\end{aligned}
$$

## Commutativity

$$
\begin{aligned}
& a+b=b+a \\
& a \times b=b \times a
\end{aligned}
$$

## Partial Distributivity

$$
a(b+c)=a b+a c
$$

When $\quad a \neq \pm \infty$ or

$$
b c>0 \text { or }
$$

$$
(b+c) / 0=\Phi
$$

## Transreal Acceptance

- Transreal arithmetic proved consistent by machine proof (2006)
- Transreal (2016) and transcomplex (2014) arithmetic proved consistent by construction from, respectively, the real and complex numbers
- So transreal and transcomplex arithmetic are consistent if real arithmetic is and can only be disproved if real arithmetic is


## Projective Geometry



## Transreal-Number Line

## $\Phi$


$\infty$

## Transcomplex Plane

Revolution of the transreal number line

## $\Phi$

## Transangle

Real and nullity angles are arc length divided by radius


But where are the infinity angles?

## Transangle

The infinity angle is the winding at the apex of the unit cone


## Transreal Acceptance

- Transreals have a geometrical model
- Transcomplexes have a geometrical model
- Transreal angles have a geometrical model


## Transtangent



## Transtangent


$\Phi$

## Transtangent

- Is defined for all transreal angles
- Is single valued everywhere
- The real values of the transtangent have period $\pi$
- The infinite values of the transtangent have period $2 \pi$
- The nullity value of the transtangent occurs at $-\infty$, $\infty, \Phi$


## Transtangent

- Transtangent definition on a triangle is consistent with
- Transtangent definition as a power series


## Transmathematics

- Transalgebra including Trans-Boolean Algebra
- Transreal \& transcomplex elementary functions
- Transreal and transcomplex calculus
- Transvectors


## Transmathematics

- Transreal arithmetic is so general it is beginning to solve problems in logic
- Transreal calculus is so general that it is beginning to solve problems in real calculus and mathematical physics (no empirical verification yet?!)


## Transreal-Number Line

$\Phi$

$\infty$

## Nullity Force

- There is no component of nullity on the extended-real-number line so nullity forces have no, that is zero, effect on the extended-real universe where we live


## Newton's Law 1

- A mass is accelerated only by a positive or negative force, not a zero or nullity force


## Newton's Law 2

- $F=m a$ when $0<m<\infty$ and $a$ is transreal
- $a=F$ / $m$ when $0<m<\infty$ and $F$ is transreal
- $m=F / a$ when $a, F$ are transreal. When the computed mass is real, it is determined. When the computed mass is nullity, the true, finite mass, is hidden (but can be discovered by its gravitational effects)


## Newton's Law 3

- To any action, F, there is always an opposite and equal reaction, -F


## Information

- Real numbers have more information than infinite numbers because both the sign and magnitude of real numbers is non-absorptive but the magnitude of infinite numbers is absorptive
- Infinite numbers have more information than nullity because nullity's sign and magnitude is absorptive
- Hypothesis: physical systems always adopt the transreal configuration with the highest possible information


## Black Hole

- Suppose we have two, same charged, massive, point particles at the singularity of a black hole
- Attraction $F_{g}=G \frac{m_{1} m_{2}}{r^{2}}=G \frac{m_{1} m_{2}}{r^{2}}=\infty$
- Repulsion $F_{e}=k_{e} \frac{q_{1} q_{2}}{r^{2}}=k_{e} \frac{q_{1} q_{2}}{r^{2}}=-\infty$
- Nett force $F=F_{g}+F_{e}=\infty-\infty=\Phi$


## Black Hole

- The particles are bound by a nullity force at the singularity so are free to move but are not compelled to move
- A quantal fluctuation in position may move some effective mass away from the singularity - if it inflates, it may leave the event horizon, if not it falls back into the singularity in a convection current


## Black Hole

- The convection current perturbs the event horizon:
- Evaporation is faster than predicted by Hawking because a bumpy and roiling event horizon has a larger surface area than a spherical one
- Heating outside the event horizon is nonmonotonic because a bumpy and roiling event horizon accelerates local particles differentially


## Electronics

- For the last one hundred years, electronic circuits have been known with properties 0/0
- The circuits are measurable
- There is no theory of what to do with them


## Serial Computation

- Maximises latency to first and successive solutions
- Efficiency decreases with increasing ratio of processor to memory speed because of memory bottleneck
- Efficiency decreases with increasing size


## Parallel Computation

- Minimises latency to first solution
- Efficiency decreases with increasing ratio of processor to memory speed
- Efficiency decreases with increasing size
- An infinitely fast or big parallel von Neuman machine has efficiency zero and does zero work


## Dataflow

- Maximises concurrency and throughput
- Processors run at memory speed
- Constant efficiency


# Serial Processing 

Instruction 1 Data 1

Instruction 2

Instruction 3

Instruction n

# Serial Processing 

Instruction 1

Instruction 2 Data 1

Instruction 3

Instruction n

# Serial Processing 

Instruction 1

Instruction 2

Instruction 3 Data 1

Instruction n

# Serial Processing 

Instruction 1

Instruction 2

Instruction 3

Instruction n Data 1

# Serial Processing 

Instruction 1 Data 2

Instruction 2

Instruction 3

Instruction n

# Serial Processing 

Instruction 1

Instruction 2 Data 2

Instruction 3

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Instruction 3 Data n

Instruction n

# Serial Processing 

Instruction 1

Instruction 2

Instruction 3

Instruction n Data n

## Slipstream Computer

- Dataflow means travel time is proportional to distance so never stalls on memory bottleneck
- Dataflow means I/O bandwidth is independent of the number of mills
- Totallity means that if a program compiles it has no logical exceptions so it can crash only on a physical fault
- Totallity means pipelines never break


# Slipstream Processing 

Instruction 1 Data 1

Instruction 2

Instruction 3

Instruction n

# Slipstream Processing 

Instruction 1 Data 2

Instruction 2 Data 1

Instruction 3

Instruction n

# Slipstream Processing 

Instruction 1 Data 3

Instruction 2 Data 2

Instruction 3 Data 1

Instruction n

# Slipstream Processing 

Instruction 1 Data n

Instruction 2 Data 3

Instruction 3 Data 2

Instruction n Data 1

# Slipstream Processing 

Instruction 1

Instruction 2 Data n

Instruction 3 Data 3

Instruction $n$ Data 2

# Slipstream Processing 

Instruction 1

Instruction 2

Instruction 3 Data n

Instruction n Data 3

# Slipstream Processing 

Instruction 1

Instruction 2

Instruction 3

## Instruction $n$ Data $n$

## FPGA Prototype



## Architectural Prototype

- Token = 12-bit header + 80-bit float datum
- 64 k mills per chip
- 2 M mills per board
- 16 M mills per cabinet
- 20 kW per unweighted Wassenaar Peta FLOP (PWFLOP)


## I/O

- Systolic arrays have one dimensional I/O which has linear scaling and is impossible to fabricate
- Architectural prototype uses zero dimensional I/O which has constant (no) scaling and can be fabricated


## Relative Addressing

- Fixed size, relative address implements an address horizon in an arbitrarily large machine and maintains constant computational efficiency regardless of the size of the machine
- Small horizon keeps the token header small


## Processor Grid

- Square grid of mills
- Pipelined communication not just nearest neighbour



## Slipstream

- A grid of mills may be arranged in any dimensionality of space (2D is convenient for chips!)
- The nodes of the grid are coloured by the configuration state of the mills
- A Turing program is a directed graph in a grid
- A slipstream program is an acyclic graph in a grid


## Slipstream

- Slipstream programs execute in a cadence (period) of the longer of the input and output times
- Programs with shared data, such as molecular dynamics, may have many copies of a program that share data so the average cadence is less than one and the limit of the cadence, with increasing machine size, can be zero!


## Slipstream

- A practical slipstream machine cannot achieve a cadence of zero
- But the ratio of the execution time of a practical slipstream machine versus a practical von Neumann serial or parallel machine can be infinity


## Transreal Acceptance

- Construction from the real numbers
- Geometrical models
- Today, machine proofs of consistency
- Tomorrow, many useful applications

